Power Efficiency and Delay Tradeoff of Energy Efficient Ethernet Protocol

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August 22, 2017



- Power Saving Problem
- System Modeling
- Performance Tradeoff
- Conclusion



## Power Saving Problem

## System Modeling

## Performance Tradeoff

## Conclusion

#### Features of Ethernet





- P. J. Winzer, "Beyond 100G Ethernet," IEEE Communications Magazine, vol. 48, pp. 26–30, July 2010.
- P. Reviriego, K. Christensen, J. Rabanillo, and J. A. Maestro, "An initial evaluation of Energy Efficient Ethernet," IEEE Communications Letters, vol. 15, pp. 578–580, May 2011.
- B. Kohl, "10GBASE-T power budget summary," 2007.

#### **EEE Protocol**

#### • Go to sleep when idle



- Sleeping strategies:
  - $\tau \& N$  policy:  $\tau$  and N are finite
  - $\tau$  policy:  $N \to \infty$
  - *N* policy:  $\tau \to \infty$



Tradeoff between power efficiency and delay
 N and τ ↑ → LPI Length ↑ → power efficiency ↑
 delay ↑

• What are the rules to select *N* and  $\tau$ ?



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#### New Feature of EEE Queuing Systems

## EEE protocol is an M/G/1 queue with vacations modulated by arrival process



 $\leftarrow$ A working cycle (*C*) $\rightarrow$ 

#### Failure of Classical Analytical Method

 Example: classical P-K formula fails to describe the mean delay of EEE protocol

$$D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{\overline{V^2}}{2\overline{V}} + \overline{X}$$



#### Failure of Classical Analytical Method

 Example: classical P-K formula fails to describe the mean delay of EEE protocol

$$D \neq \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{\overline{V^2}}{2\overline{V}} + \overline{X}$$







#### Arrival Event Tree of Vacation Periods

#### Six mutually-exclusive events



•  $h_n = \Pr\{n \text{ arrivals during a vacation period } V\}$   $H(z) = \sum_{n=0}^{\infty} h_n z^n$   $= e^{-\lambda T_w(1-z)}$  $\times \left[\sum_{n=0}^{N-1} e^{-\lambda T_s} \frac{(\lambda T_s)^n}{n!} (z^N - z^n) - \sum_{n=0}^{N-2} e^{-\lambda \tau} \frac{(\lambda \tau)^n}{n!} (z^N - z^{n+1})\right]$ 

## Mean Vacation Time and Mean Cycle Time

- H'(1): the mean number of arrivals during vacation
- By Little's Law, mean vacation time  $\overline{V}$  $\overline{V} = \frac{H'(1)}{\lambda}$
- Mean cycle time

$$\bar{C} = \frac{\bar{V}}{1-\rho}$$

#### Power Efficiency $\eta$

 $\eta = \frac{\text{average power saved in one cycle by an LPI}}{\text{averge power of one cycle if EEE is not used}}$  $= \frac{(\overline{V} - T_S - T_W) \times (\varphi_h - \varphi_l)}{\overline{C} \times \varphi_h}$  $= \left(1 - \frac{T_W + T_S}{\overline{V}}\right) \cdot \frac{(1 - \rho) \times (\varphi_h - \varphi_l)}{\varphi_h}$ 



#### Delay Analysis



• D. P. Bertsekas, R. G. Gallager, and P. Humblet, *Data networks*, vol. 2. Prentice-Hall International New Jersey, 1992.

#### Mean Residual Time

$$R = E[R_i | \xi = 0] \times Pr\{\xi = 0\} + E[R_i | \xi = 1] \times Pr\{\xi = 1\}$$
$$= E[R_i | \xi = 0] \times (1 - \rho) + E[R_i | \xi = 1] \times \rho$$
$$\xi = \begin{cases} 0, & \text{if frame arrives during a vacation period} \\ 1, & \text{if frame arrives during a busy period} \end{cases}$$

$$E[R_i|\xi = 1] = \frac{1}{2\rho}\lambda \overline{X^2}$$
: independent of arrival process  
 $E[R_i|\xi = 0] = ?$ : dependent on arrival process

• D. P. Bertsekas, R. G. Gallager, and P. Humblet, *Data networks*, vol. 2. Prentice-Hall International New Jersey, 1992.

#### Residual Vacation Time of Each Arrival

## • Given $V_n$ , number of arrivals during residual vacation time seen by a frame is determined



 $E[R_i|\xi = 0] = \sum_{n=1}^{\infty} E[R_i|\xi = 0, \text{ frame } i \text{ arrives in a } V_n] \cdot P_n$ •  $P_n$ : conditional probability that a frame arrives in a  $V_n$ .

• Applying Little's Law  $\lambda E[R_i|\xi = 0] = \sum_{n=1}^{\infty} \lambda E[R_i|\xi = 0, \text{ frame } i \text{ arrives in a } V_n] \cdot P_n$   $= \sum_{n=1}^{\infty} E[Q_i|\xi = 0, \text{ frame } i \text{ arrives in a } V_n] \cdot P_n$   $= \sum_{n=1}^{\infty} \left[\frac{(n-1)+(n-2)+\dots+1+0}{n}\right] \cdot \frac{n \cdot h_n}{H'(1)}$   $= \frac{H''(1)}{2H'(1)}.$  Theorem 1: The mean delay of EEE systems is given by:

$$D = \frac{\lambda X^2}{2(1-\rho)} + \frac{H''(1)}{2\lambda H'(1)} + \bar{X}$$
  
• Classical P-K Formula  

$$D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{\overline{V^2}}{2\overline{V}} + \bar{X}$$



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Timer vs. Counter

# $\eta_{\tau \& N} \approx \eta_{\tau} \approx \eta_{N}$ $D_{\tau \& N} = \min\{D_{\tau}, D_{N}\}$



#### Explanation: Wakeup Triggers of $\tau \& N$ policy

•  $\tau \& N$  policy can adaptively use  $\tau$  and N to wake up according to instantaneous arrival rate if  $\tau = \frac{N-1}{\lambda}$ 



For a given steady state traffic rate  $\lambda$ , the selection of parameters  $\tau$  and *N* should comply with the following condition:

$$\frac{N-1}{\tau} = \lambda.$$

#### Power Efficiency versus Mean Delay

 Excessively large τ and N degrade delay performance while marginally enhancing the power efficiency



Based on EEE 1, parameter *N* can be selected according to a given average delay requirement *D* from the expression of  $D_N$ 

$$D_N \approx \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{(N+\lambda T_w)^2 - N}{2\lambda(N+\lambda T_w)} + \overline{X}$$

#### Conclusions

- Develop a new approach to analyze the *M/G/1* queue with vacations governed by the arrival process
- Derive a generalized P-K formula of mean delay
- Provide two rules to select appropriate  $\tau$  and N



## Thank You!