

# Modular Optical Cross-Connects (OXC) for Large-Scale Optical Networks

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**Abstract**—Due to the explosive growth of traffic demands, large-scale optical cross-connects (OXC) are highly desired in next-generation optical networks. However, the scalability of classical OXC is not good in terms of the size of optical switching modules and cabling complexity to meet application requirement. To address this issue, this letter proposes a two-phase approach to construct modular large-scale OXC, using a set of small-size OXC modules. We first decompose each optical space switch (OS) in a classical OXC to a Clos network and then merge the corresponding links and decomposed optical switching modules of different Clos networks to form small-size OXC modules, which finally yields a modular OXC. We show that the proposed modular OXC has good scalability and acceptable physical-layer performance.

**Index Terms**—Modularity, optical cross-connect (OXC), optical network.

## I. INTRODUCTION

**O**PTICAL cross-connect (OXC) based switching node is one of essential components in wavelength routing optical networks [1], [2]. With the explosive growth of traffic demands driven by high-bandwidth applications, the number of wavelength-division-multiplexing (WDM) fiber links between two adjacent switching nodes in the networks are being forced to increase fast due to the limitation of the capacity per link [2]–[4]. Therefore, large-scale OXC are highly desired.

However, the scalability of classical OXC is poor in terms of the size of each module and the number of fiber links. As Fig. 1 shows, an  $N \times N$  OXC with  $w$  wavelength channels consists of  $w N \times N$  optical space switches (OSs) sandwiched by  $N$   $1 \times w$  input wavelength demultiplexers (DeMux) on the left and  $N$   $w \times 1$  wavelength multiplexers (Mux) on the right. The OS is typically a micro-electro mechanical system (MEMS). Though free-space-optics based 3-D MEMSs with port count larger than 300 have been commercial, such kind of MEMSs is somewhat bulky for practical applications [5]–[7]. On the other hand, cabling complexity has been an important concern in the application of large-scale switching node [8]. In a classical OXC, there is one and only one fiber link

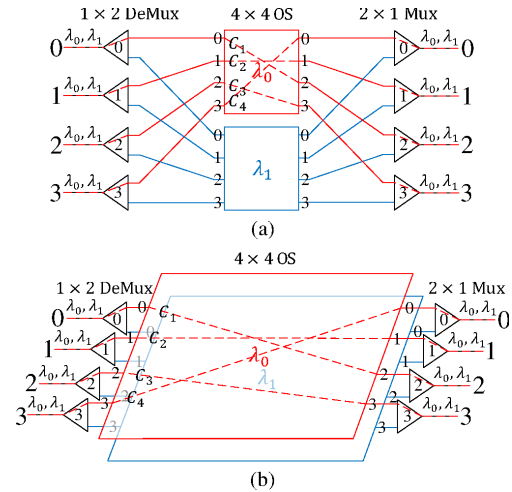


Fig. 1. A  $4 \times 4$  OXC: (a) planar structure, and (b) stereo representation.

between the optical components in two adjacent stages, and thus the number of fibers is  $2Nw$ . If we want to construct, for example, a  $320 \times 320$  OXC with 24 wavelengths directly from 24  $320 \times 320$  MEMSs, the number of fiber links will be 15360, which will impose a heavy maintenance burden on the network operators. Thus, as Ref. [7] pointed out, the construction of the OXC from large-scale MEMSs is not practical.

To address scalability issue, several kinds of large-scale OXC have been proposed. Ref. [3] constructed a large-scale two-stage-routing OXC using a set of small-size wavelength selective switches (WSSs). However, such OXC cannot freely switch a wavelength from any input to any output due to the use of coarser-granular switching devices. In [4], a modular OXC was constructed from a set of small-scale OXC connected in a ring topology. Such kind of OXC has two drawbacks. Firstly, its network structure is completely different from that of the classical OXC and thus requires a different routing algorithm. This means the operation principle is not compatible with the classical OXC. Secondly, the insertion loss of the OXC remarkably increases with the port count, which clearly limits the scalability.

In this letter, we propose a two-phase approach to construct a modular OXC, which is an interconnection network of small-size OXC modules. We start with a large-scale classical OXC. In phase 1, we decompose each OS in the OXC into a three-stage Clos network to modularize the OSs, which however doubles the number of interconnection links in the OXC. In phase 2, we merge the corresponding links and decomposed switching modules of different Clos networks to

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form a small-scale OXC module, such that the number of fiber links that interconnect the small-size OXC modules is reduced to  $4N$ . After such procedure, we transform a large-scale classical OXC to a modular OXC, which possesses the following advantages:

- (1) It can be constructed from a set of small-size OXC modules, each of which is easy to synthesize and may even be integrated in a chip [9] in the near future;
- (2) The number of fiber links that interconnect the OXC modules is only  $2/w$  of that of the traditional OXC, which remarkably reduces the cabling complexity;
- (3) It is non-blocking at each wavelength, and its routing algorithm is compatible with that for classical Clos networks;
- (4) Its insertion loss is acceptable, and does not increase fast with the port count.

## II. CONSTRUCTION APPROACH FOR MODULAR OXCS

To facilitate the presentation, we denote an  $N \times N$  OXC, as  $\mathcal{O}_A(N, w)$ . For example, the OXC in Fig. 1(a) is an  $\mathcal{O}_A(4, 2)$ .

We number the DeMuxes by  $0, 1, \dots, N-1$ , and the Muxes by  $0, 1, \dots, N-1$  from top to bottom. Each DeMux (or Mux) carries a wavelength set  $\Lambda = \{\lambda_0, \lambda_1, \dots, \lambda_{w-1}\}$ . Similarly, we label the OSs by  $0, 1, \dots, w-1$ . The  $l$ th OS performs switching for the signals at wavelength  $\lambda_l$ , where  $l = 0, 1, \dots, w-1$ .

We denote a connection as  $C(a, \beta, \lambda_l)$ , if it is a connection from wavelength  $\lambda_l$  at input  $a$  to that at output  $\beta$ , where  $a, \beta = 0, 1, \dots, N-1$ . The procedure that the OXC sets up connection  $C(a, \beta, \lambda_l)$  is as follows. At input  $a$ , the DeMux filters out  $\lambda_l$  from the input WDM signal and sends  $\lambda_l$  via its  $l$ th output to the  $a$ th input of OS  $l$ . Through reconfiguration, OS  $l$  can switch  $\lambda_l$  to the  $l$ th input of Mux  $\beta$  via its  $\beta$ th output. An example is  $C_1(0, 2, \lambda_0)$  in Fig. 1.

This letter also refers to an OS in the classical OXC as a wavelength switching plane or a wavelength plane. To illustrate this concept, we represent the OXC in a stereo form, in which the DeMuxes/Muxes are respectively laid out from back to front, and the wavelength planes are stacked in parallel from top to bottom. As an example, Fig. 1(b) plots the stereo presentation of  $\mathcal{O}_A(4, 2)$  in Fig. 1(a).

When  $N$  becomes large, the port count of each OS is large, and the scalability of classical OXCs is poor. In this section, we show how a large-scale OXC can be constructed from a set of small-scale OXC modules.

### A. Decomposition

As Ref. [10] demonstrates, decomposition is a typical way for network modularization. We thus perform OS decomposition in the first step, to scale down the size of OSs in the OXC when  $N$  is large. In particular, we decompose each OS to a symmetric  $N \times N$  three-stage Clos network, where there are:

- (1)  $r \ n \times n$  OSs in the first stage,
- (2)  $n \ r \times r$  OSs in the second stage,
- (3)  $r \ n \times n$  OSs in the third stage, and
- (4) exactly one fiber link between the OSs at two adjacent stages, where  $N = rn$ .

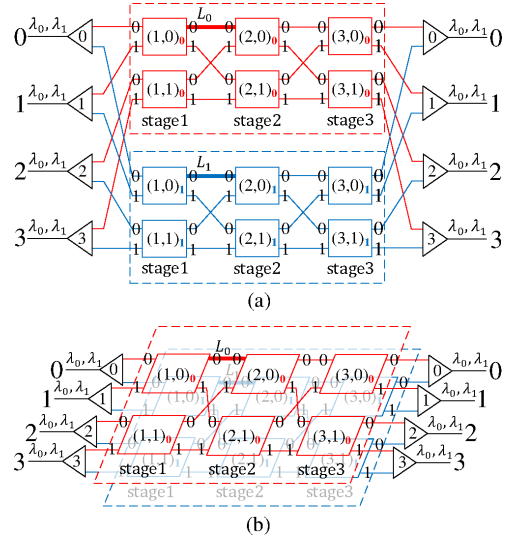


Fig. 2. A  $4 \times 4$  network  $\mathcal{O}_B(4, 2)$  after decomposition: (a) planar structure, and (b) stereo representation.

It is clear that each wavelength plane after the decomposition preserves the non-blocking property, according to the rearrangeably non-blocking condition of Clos networks.

After that, we obtain a switching network called  $\mathcal{O}_B(N, w)$ . For example, OXC  $\mathcal{O}_A(4, 2)$  in Fig. 1(a) changes to  $\mathcal{O}_B(4, 2)$  in Fig. 2(a), after each  $4 \times 4$  OS or wavelength plane in Fig. 1(a) is decomposed to a  $4 \times 4$  Clos network, in which  $r = n = 2$ . Accordingly, Fig. 2(b) is the stereo representation of Fig. 2(a).

Though the decomposition phase scales down the port count of each OS, it introduces  $2Nw$  fiber links in additional, such that the number of fiber links that interconnect the network components is twice that of the original OXC. We thus need the second phase to reduce the number of fiber links.

### B. Merge

At the  $l$ th wavelength plane, we denote an OS in  $\mathcal{O}_B(N, w)$  as  $(1, i)_l$  if it is the  $i$ th OS in the first stage, an OS as  $(2, j)_l$  if it is the  $j$ th OS in the second stage, and an OS as  $(3, k)_l$  if it is the  $k$ th OS in the third stage, where  $i, k = 0, 1, \dots, r-1$  and  $j = 0, 1, \dots, n-1$ . For example,  $(1, 1)_0$  is the 1st OS in the 1st stage at wavelength plan  $\lambda_0$ . Based on such numbering scheme, we can find that the stereo representation of  $\mathcal{O}_B(N, w)$  has the properties as follows:

- P1: OSs  $(1, i)_0, (1, i)_1, \dots, (1, i)_{w-1}$  are vertically aligned  
OSs  $(2, j)_0, (2, j)_1, \dots, (2, j)_{w-1}$  are vertically aligned  
OSs  $(3, k)_0, (3, k)_1, \dots, (3, k)_{w-1}$  are vertically aligned;
- P2: Each  $1 \times w$  DeMux at the input side connects with a set of  $w$  vertically-aligned OSs at the first stage;
- P3: Each  $w \times 1$  Mux at the output side connects with a set of  $w$  vertically-aligned OSs at the third stage.

These properties hold since all the Clos networks at different wavelength planes have the same topology. For example, OS  $(1, 0)_0$  and OS  $(1, 0)_1$  in Fig. 2(b), are vertically aligned. This is also true for OS  $(2, 0)_0$  and OS  $(2, 0)_1$  in Fig. 2(b).



Fig. 3. Two vertically-aligned links can be implemented by a multi-fiber link containing two fibers.

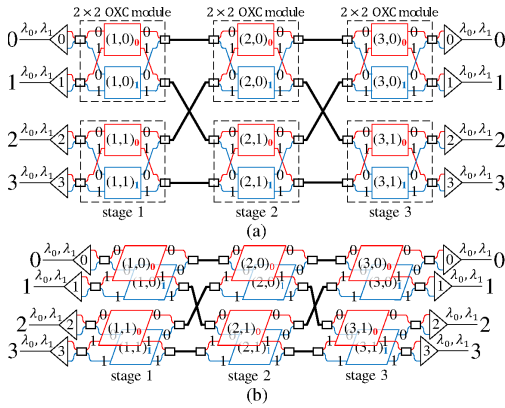


Fig. 4. A  $4 \times 4$  modular OXC  $\mathcal{O}_C(4, 2)$ : (a) planar structure, and (b) stereo representation.

Furthermore,  $1 \times 2$  DeMux 3 connects with OSs  $(1, 1)_0$  and  $(1, 1)_1$ , and  $2 \times 1$  Mux 1 connects with OSs  $(3, 0)_0$  and  $(3, 0)_1$ . It follows that  $w$  links between two sets of vertically-aligned OSs are also vertically aligned. For example, in Fig. 2, link  $L_0$  between OS  $(1, 0)_0$  and OS  $(2, 0)_0$  is vertically aligned with link  $L_1$  between OS  $(1, 0)_1$  and OS  $(2, 0)_1$ .

Each set of  $w$  vertically-aligned links in  $\mathcal{O}_B(N, w)$  can be implemented by a multi-fiber cable containing  $w$  fibers between two Multi-fiber PushOn (MPO) connectors. As Fig. 3 illustrates, the two links  $L_0$  and  $L_1$  in Fig. 2 can be implemented by a multi-fiber cable containing two fibers between a pair of MPO connectors. Note that, a currently commercial MPO connector only contains up to 24 fibers. Thus, if  $w > 24$ , we need a bound of multi-fiber cables to replace a group of  $w$  vertically-aligned links. We expect that the number of fibers contained in an MPO connector could be large in the future.

We thus replace each set of  $w$  vertically-aligned links in  $\mathcal{O}_B(N, w)$  by a multi-fiber cable, resulting in a new switching network called  $\mathcal{O}_C(N, w)$ . For example,  $\mathcal{O}_C(4, 2)$  in Fig. 4 is obtained by replacing each set of two vertically-aligned links in Fig. 2 by a multi-fiber cable.

According to P1 and P2, each set of  $w$  vertically-aligned OSs at the first stage is now sandwiched by  $n$  MPO connectors on the left and  $n$  MPO connectors on the right, which forms an  $n \times n$  OXC module. In the same way, we obtain  $n$   $r \times r$  OXC modules at the second stage and  $r$   $n \times n$  modules at the third stage. We thus refer to  $\mathcal{O}_C(N, w)$  as  $N \times N$  modular OXC.

As Fig. 4 shows, the total number of links that interconnect the small-size OXC modules in modular OXC  $\mathcal{O}_C(N, w)$  is  $4N$ , which is only  $2/w$  of the number of the links that interconnect the OSs in the traditional OXC. Since the links inside each OXC module are encapsulated in a standard module, they do not increase the cabling complexity of the whole system. We thus claim that, compared with the classical OXC, the modular OXC reduces the cabling complexity by a fact of  $w/2$ . For example, we only need 1280 multi-fiber cables to construct a  $320 \times 320$  modular OXC  $\mathcal{O}_C(320, 24)$ .

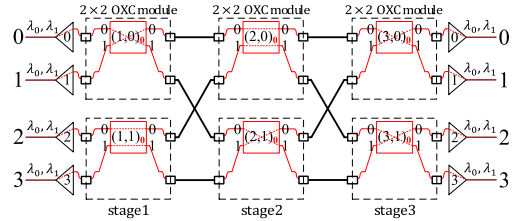


Fig. 5. Route assignments in the wavelength plane  $\lambda_0$  of a modular OXC.

Also, there are  $N$  DeMuxs,  $N$  Muxs,  $6N$  MPO connectors,  $2rw$   $n \times n$  OSs, and  $nw$   $r \times r$  OSs in a modular OXC  $\mathcal{O}_C(N, w)$ . As an example,  $\mathcal{O}_C(320, 24)$  contains 320 DeMuxs, 320 Muxs, 1920 MPO connectors, 960  $16 \times 16$  OSs, and 384  $20 \times 20$  OSs.

### III. ROUTING PROCESS OF MODULAR OXCS

In this section, we show how to assign routes for connection requests in modular OXC  $\mathcal{O}_C(N, w)$ . As we know, there is no wavelength converter in the modular OXC. It follows that, once the wavelength of a connection request is given, it can only pass through the OSs in the corresponding wavelength plane. Thus, the routing assignment of requests at different wavelengths can be coped with independently.

Suppose there are  $m$  requests  $C_1, C_2, \dots, C_m$  at wavelength  $\lambda_l$ , where  $m \leq N$ . We first ignore all the OSs at other wavelength planes. Clearly, the OSs at wavelength  $\lambda_l$  constitutes a Clos network, as Section II shows. For example, we want to assign routes for  $C_1(0, 2, \lambda_0)$ ,  $C_2(1, 1, \lambda_0)$ ,  $C_3(2, 3, \lambda_0)$ , and  $C_4(3, 0, \lambda_0)$ . We ignore the OSs at wavelength  $\lambda_1$  and obtain a Clos network in Fig. 5.

The core issue of the routing in Clos network is to assign the OSs at the second stage to the requests. The routing constraint at each wavelength is that two requests cannot pass through the same OS at the second stage if they share the same OS at the first stage or the same OS at the third stage. Currently, there are a lot of classical algorithms available to solve the routing problem in Clos networks. For example, we can use the Karol's algorithm in [11]. After such routing process, we can assign routes for  $C_1$  through  $C_4$ , as Fig. 5 shows.

### IV. PHYSICAL-LAYER PERFORMANCE

In this section, we carry out a proof-of-concept experiment to demonstrate that the proposed modular OXC is feasible in terms of physical-layer performance. We show that the end-to-end insertion loss of each connection is acceptable.

In our experiment, we take the  $4 \times 4$  modular OXC in Fig. 4 as an example. We consider two connections C1 where all the  $2 \times 2$  OSs are in bar state, and C2 where all the OSs are in cross state, to demonstrate the transmission performance difference of different connections.

Fig.6 plots the experimental setup. We first produce a 6-dBm 10-Gb/s optical signal via a 10-G directly modulated laser (DML) operating at 1550.54 nm, which is driven by a 10-Gb/s non-return-to-zero on-off-keying (NRZ-OOK) electrical signal generated by a pulse pattern generator (PPG, Anritsu MP1800A). We feed this optical signal to an optical connection, containing two  $8 \times 8$  arrayed waveguide gratings (AWGs) and three  $2 \times 2$  magneto-optical switches

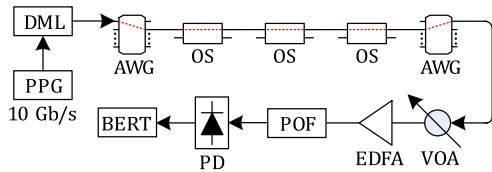


Fig. 6. Experimental setup.

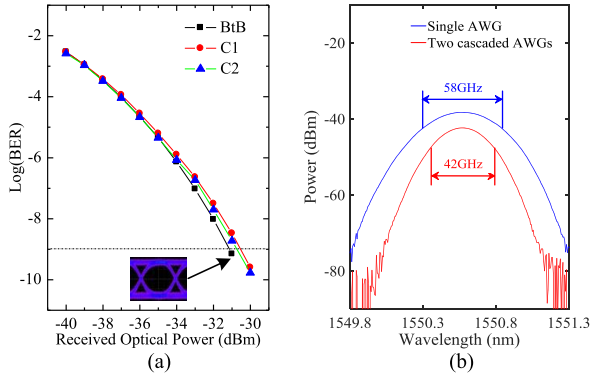


Fig. 7. Experimental results: (a) BER performance, and (b) pass band of two cascaded AWGs.

(Primanex MagLight<sup>TM</sup>). Herein, we only use one input and two outputs of each  $8 \times 8$  AWG to fulfill the function of a  $1 \times 2$  DeMux (or  $2 \times 1$  Mux). After the optical signal passes through C1 (or C2), the received optical power is  $-4.2$  dBm (or  $-3.8$  dBm). Thus, the insertion loss suffered by C1 (or C2) is 10.2 dB (or 9.8 dB).

At the receiver side, the optical signal in turn passes through a variable optical attenuator (VOA) for received optical power adjustment, an erbium-doped fiber amplifier (EDFA) for power amplification, and a programmable optical filter (POF, Finisar Waveshaper) for amplified spontaneous emission noise elimination. The electrical signal is retrieved from the optical signal by a photodiode (PD), and finally captured by a bit-error rate tester (BERT) to measure the bit-error-ratio (BER). In the transmission performance measurement, we consider the following two conditions:

- (1) Back-to-back (BtB): the optical signal passes through a fiber link;
- (2) C1/C2: the optical signal goes through C1/C2.

Experimental results are plotted in Fig. 7. It can be seen from Fig. 7(a) that the BER curves of the NRZ-OOK signals in C1 and C2 are almost the same, which demonstrates that the transmission performance difference of various connections is ignorable if the magneto-optical switch is used. Fig. 7(a) also shows that the power penalty of C1 is only 0.6 dB compared to BtB at  $\text{BER} = 10^{-9}$ , which is mainly attributed to the fact that the central frequency of the optical signal is drifted by the DML chirp, and the optical signal suffers from filtering effect [12] when it passes through two AWGs. As Fig. 7 (b) shows, the pass band of a single 100-GHz AWG is 58 GHz, and that of two cascaded 100-GHz AWG reduces to 42 GHz. From the result in Fig. 7, we conclude that the proposed modular OXC is feasible in terms of physical-layer performance.

The proposed  $N \times N$  modular OXC has a good scalability in terms of optical insertion loss even when the port count  $N$

is large. In the modular OXC, each connection passes through 1 DeMux, 6 MPO connectors, 3 OSs, and 1 Mux. Compared with that in the classical OXC, each connection in the modular OXC traverses 6 more MPO connectors and 2 more OSs. The maximum insertion losses of a 100-GHz AWG-based Mux (or DeMux), an MPO connector, and a 32-port OS are 5.5 dB [13], 0.35 dB, and 1.4 dB [14], respectively. Therefore, the worst-case insertion loss of a connection in the modular OXC is 17.3 dB, about 4.9 dB more than that in the classical OXC. If the modular OXC is deployed in optical networks, an EDFA will be needed at each output, which may be the cost of modularization. Even so, the modular OXC has an advantage that its insertion loss does not increase fast with the port count.

## V. CONCLUSION

This letter proposes an approach to design modular OXCs, which possess the following advantages. First, it is constructed from a set of small-scale OXC modules, each of which is easy to synthesize and may even be integrated in a chip in the near future. Second, the number of fiber links that interconnect the OXC modules is only  $2/w$  of that of the classical OXC. Third, it is non-blocking at each wavelength, and the routing algorithm is compatible with that for classical Clos networks. Furthermore, its insertion loss is acceptable, and does not increase fast with the port count.

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