

Power Efficiency and Delay Tradeoff of Energy Efficient Ethernet Protocol

Xiaodan Pan, Tong Ye, Tony T. Lee, Weisheng Hu pxd0506@sjtu.edu.cn

State Key Lab of Advanced Optical Communications and Networks Shanghai Jiao Tong University

Background

- **Energy Efficient Ethernet Protocol**
- **Nacation Model**
- **Power Efficiency**
- **P-K Formula of Mean Delay**
- **Tradeoff and Parameter Selections**
- Conclusion

Widely Applied Ethernets

LAN WAN^[3]

[1] A. Greenberg, J.R. Hamilton, N. Jain, et al, "VL2: a scalable and flexible data center network," *Proc. ACM SIGCOMM*, 2009, pp. 51-62. [2] M. Huynh, P. Mohapatra, "Metropolitan Ethernet Network: A move from LAN to MAN," *Computer Networks*, vol. 51, pp. 4867-4894, Dec 2007. [3] A. Kasim, P. Adhikari, N. Chen, et al, "Ethernet: From LAN to the WAN," in *Delivering Carrier Ethernet*, 1st ed., New York: McGraw-Hill, 2007, pp.3-43.

Growing of Ethernet Devices

The number of devices is huge and still grows rapidly.

Network Equipment Market Scale and Forecast of China

[1] http://www.ccwresearch.com.cn/view_point_detail.htm?id=557063

[2] R. Bolla, R. Bruschi, F. Davoli, and F. Cucchietti, "Energy efficiency in the future Internet: A survey of existing approaches and trends in energyaware fixed network infrastructures," *IEEE Communications Surveys Tutorials*, vol. 13, pp. 223–244, Second 2011.

Increase of Data Rate

[1] P. J. Winzer, "Beyond 100G Ethernet," *IEEE Communications Magazine*, vol. 48, pp. 26–30, July 2010. [2] P. Reviriego, K. Christensen, J. Rabanillo, and J. A. Maestro, "An initial evaluation of Energy Efficient Ethernet," *IEEE Communications Letters*, vol. 15, pp. 578–580, May 2011. [3] B. Kohl, "10GBASE-T power budget summary," 2007.

Link speed (Mb/s)

Idea of Energy Saving

 IEEE 802.3az: Shut down some component during idle periods and make the system more energy proportional to load

Trace from LBNL: File Server with 1G Ethernet Link

- **Background**
- **Energy Efficient Ethernet Protocol**
- **Nacation Model**
- **Power Efficiency**
- **P-K Formula of Mean Delay**
- **Tradeoff and Parameter Selections**
- Conclusion

Energy Efficiency Ethernet Protocol

- Sleep: transition time from Active to LPI
- Wakeup: transition time from LPI to Active
- LPI: low power idle mode
- Active: packets transmission period

A Typical State Transition and Power Consumption of EEE Protocol

- \blacksquare Counter N
	- Bound the backlogged queue length
- **Timer** τ ($\tau > T_s$)
	- **Bound the delay**

- \blacksquare Counter N
	- Bound the backlogged queue length
- **Timer** τ ($\tau > T_s$)
	- **Bound the delay**

- \blacksquare Counter N
	- Bound the backlogged queue length
- **Timer** τ ($\tau > T_s$)
	- **Bound the delay**

 τ &N policy

- \blacksquare Counter N
	- Bound the backlogged queue length
- **Timer** τ ($\tau > T_s$)
	- **Bound the delay**

 τ &N policy

- \blacksquare Counter N
	- Bound the backlogged queue length
- **Timer** τ ($\tau > T_s$)
	- **Bound the delay**

τ policy and N policy

Performance Tradeoff

- Power efficiency is improved at the expanse of delay.
- How to select N and τ to optimize system performances?

Our Works

- Model BTR strategy as an M/G/1 queue with vacation time which is governed by the arrival process.
- Derive the P-K formula of mean delay.
- Demonstrate the impacts of counter and timer on performances and provide two rules to select appropriate parameters N and τ .

- **Background**
- **Energy Efficient Ethernet Protocol**
- **Nodel** Vacation Model
- **Power Efficiency**
- **P-K Formula of Mean Delay**
- **Tradeoff and Parameter Selections**
- Conclusion

Cycles of EEE Working Process

A renewal cycle $(C)=V$ acation period $(V)+B$ usy period (B)

EEE Working Process

Cycles of EEE Working Process

A renewal cycle $(C)=V$ acation period $(V)+B$ usy period (B)

EEE Working Process

■ Vacation period depends on the probability h_n $h_n = Pr\{n \text{ arrivals during a vacation period } V\}$

■ Vacation period depends on the probability h_n $h_n = Pr\{n \text{ arrivals during a vacation period } V\}$

Six mutually independent events

$a_n = Pr\{n \text{ arrivals in the interval } T_s\}$	$\frac{m_{\text{initial}}}{n!}$	
$= e^{-\lambda T_s (\frac{\lambda T_s}{n!})}$	Conuter expires	
Counter expires first	Comther expires during T_s	Even 1
$\frac{m_{\text{initial}}}{n}$	Now $\frac{m_{\text{initial}}}{n}$	
Counter expires first	Comther expires during LP[$(n = N)$	Event 3

\n**First arrival**

\nInterval I_0

\n $I_0 \geq T_s$

\nCounter expires first $(n = 1, 2, ..., N-1)$

\nEvent 5

\nEvent 6

\nTime expires first $(n = 1, 2, ..., N-1)$

\nEvent 6

\n**Hint**

\n $\frac{I_0}{V_{\text{in}} + V_{\text{in}} + V_{\text{in}} + V_{\text{in}} + V_{\text{in}} - V_{\text{out}}$

\n**Hint**

\n**Hint**

\n**Hint**

\n**Hint**

\n**Hint**

\n**Hint**

\n

 a_n (0 < n < N)

$a_N = 1 - \sum_{n=1}^{N-1} e^{-\lambda \tau} \frac{(\lambda \tau)^{n-1}}{(n-1)!} - \sum_{n=N+1}^{\infty} e^{-\lambda \tau} \frac{(\lambda \tau)^n}{n!}$	Number of course, we have				
$a_N = 1 - \sum_{n=1}^{N-1} e^{-\lambda \tau} \frac{(\lambda \tau)^{n-1}}{(n-1)!} - \sum_{n=N+1}^{\infty} e^{-\lambda \tau} \frac{(\lambda \tau)^n}{n!}$					
$\left.\begin{array}{c}\n\text{Counter expires first} \\ \text{during } T_s\n\end{array}\right\}$ \n	\n $\left.\begin{array}{c}\n\text{Counter expires first} \\ \text{during } T_s\n\end{array}\right\}$ \n	\n $\left.\begin{array}{c}\n\text{Counter expires first} \\ \text{during } T_s\n\end{array}\right\}$ \n	\n $\left.\begin{array}{c}\n\text{Number expires first} \\ \text{during } T_s\n\end{array}\right\}$ \n	\n $\left.\begin{array}{c}\n\text{Number expires first} \\ \text{using LPI}(n=N) \\ \text{String T}_s\n\end{array}\right\}$ \n	\n $\left.\begin{array}{c}\n\text{Number expires first} \\ \text{using LPI}(n=N) \\ \text{String T}_s\n\end{array}\right\}$ \n

П

Probability h_n

$$
a_n = \begin{cases} 0, & n = 0 \\ e^{-\lambda \tau} \frac{(\lambda \tau)^{n-1}}{(n-1)!}, & n = 1, 2, \cdots, N-1 \\ \sum_{n=0}^{N} e^{-\lambda T_s} \frac{(\lambda T_s)^n}{n!} - \sum_{n=1}^{N-1} e^{-\lambda \tau} \frac{(\lambda \tau)^{n-1}}{(n-1)!}, & n = N \\ e^{-\lambda T_s} \frac{(\lambda T_s)^n}{n!}, & n = N+1, N+2, \cdots \end{cases}
$$

 $b_n = Pr\{n \text{ arrivals during } T_w\} = e^{-\lambda T_w} \frac{(\lambda T_w)^n}{n!}$ $n!$

 $\rightarrow h_n = \sum_{k=0}^n a_{n-k} b_k \rightarrow H(z) = A(z)B(z)$

- **Mean number of arrivals during vacation is** $\overline{\alpha} = H'(1)$
	- $H(z) = \sum_{n=0}^{\infty} h_n z^n$ $n=0$
	- $\overline{\alpha}$: mean number of arrivals during vacation period
- By Little's Law, the mean vacation time \overline{V} is

$$
\bar{\alpha} = \lambda \bar{V} \rightarrow \bar{V} = \frac{\bar{\alpha}}{\lambda}
$$

- **Background**
- **Energy Efficient Ethernet Protocol**
- **Nacation Model**
- **Power Efficiency**
- **P-K Formula of Mean Delay**
- **Tradeoff and Parameter Selections**
- Conclusion

Power Efficiency η

$$
\eta = \frac{\varphi_h - \varphi_{EEE}}{\varphi_h} = \left(1 - \frac{T_w + T_s}{\overline{V}}\right) \cdot \frac{(1 - \rho) \times (\varphi_h - \varphi_l)}{\varphi_h}
$$

When N and τ come to infinite

$$
\eta^* = \lim_{\overline{V} \to \infty} \eta = \frac{(1-\rho) \times (\varphi_h - \varphi_l)}{\varphi_h}
$$

States

- **Background**
- **Energy Efficient Ethernet Protocol**
- **Nacation Model**
- **Power Efficiency**
- **P-K Formula of Mean Delay**
- **Tradeoff and Parameter Selections**
- Conclusion

Classical M/G/1 with Vacation System

 Vacation time distribution is independent of the arrival process $H(z) = \sum_{n=0}^{\infty} h_n z^n = \sum_{n=0}^{\infty} \int_{0}^{\infty} (\lambda x)^n \frac{1}{n!}$ \boldsymbol{e} ∞ $\left(\frac{1}{2}x\right)$ $\frac{1}{2}e^{-\lambda x}$ 0 $\sum_{n=0}^{\infty} \int_{0}^{\infty} (\lambda x)^{n} \frac{1}{n!} e^{-\lambda x} dV(x) z^{n}$ $n=0$ $= V^* (\lambda - \lambda z)$

Classical M/G/1 with Vacation

Failure of $H(z) = V^*(\lambda - \lambda z)$

 In EEE protocol, the vacation time is completely governed by the arrival process. Take τ policy for example

$$
V = I_0 + \tau + T_w
$$

$$
\nu(x) = \lambda e^{-\lambda(x - \tau - T_W)} (x \ge \tau + T_W) \to V^*(s) = \frac{\lambda}{\lambda + s} e^{-s(\tau + T_W)}
$$

$$
h_n = e^{-\lambda(\tau + T_w)\frac{[\lambda(\tau + T_w)]^{n-1}}{(n-1)!}} (n \ge 1) \to H(z) = ze^{-\lambda(1-z)(\tau + T_w)}
$$

• Obviously, $H(z) \neq V^*(\lambda - \lambda z)$

■ P-K Formula in the classical M/G/1 queue with vacation time

$$
D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{\overline{V^2}}{2\overline{V}} + \overline{X}
$$

For the same reason, it fails in EEE systems

The waiting time for a frame is constituted by three parts.

: mean residual time

Mean Delay

$R = E[R_i | \xi = 0] \times Pr\{\xi = 0\} + E[R_i | \xi = 1] \times Pr\{\xi = 1\}$ $= E[R_i|\xi = 0] \times (1-\rho) + E[R_i|\xi = 1] \times \rho$

 $\xi = \{$ 0, if a arrival comes during a vacation period 1, if a arrival comes during a busy period

$$
E[R_i|\xi=1] = \frac{1}{2\rho} \lambda \overline{X^{2}}^{[1]}
$$

$$
E[R_i|\xi=0] = ?
$$

Residual Vacation Time of Each Arrival

When given V_n , # of arrival during the residual vacation time seen by a frame is determined.

 $E[R_i|\xi=0]=\sum_{n=1}^{\infty}E[R_i|\xi=0, frame\ i\ arrives\ in\ a\ V_n]\cdot P_n$ $n=1$

Applying Little's Law

 $\lambda E[R_i|\xi=0]=\sum_{n=1}^{\infty}\lambda E[R_i|\xi=0, frame\;i\;arrives\;in\;a\;V_n]\cdot P_n$ $n=1$

> $=\sum_{n=1}^{\infty}E[Q_i|\xi=0, frame\ i\ arrives\ V_n]\cdot P_n$ $n=1$

$$
= \sum_{n=1}^{\infty} \left[\frac{(n-1)+(n-2)+\dots+1+0}{n} \right] \cdot \frac{n \cdot h_n}{H'(1)} = \frac{H''(1)}{2H'(1)}.
$$

Theorem 1: The mean delay of EEE systems is given by:

$$
D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{H''(1)}{2\lambda H'(1)} + \overline{X}
$$

Classical P-K Formula

$$
D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{\overline{V^2}}{2\overline{V}} + \overline{X}
$$

• When $H(z) = V^*(\lambda - \lambda z)$ holds, $D = \frac{\lambda \overline{X^2}}{2(1-\lambda)z}$ $2(1-\rho)$ + $H''(1)$ $2\lambda H'(1)$ $+ \bar{X}$ degenerate into the classical form.

- **Background**
- **Energy Efficient Ethernet Protocol**
- **Nacation Model**
- **Power Efficiency**
- **P-K Formula of Mean Delay**
- **Tradeoff and Parameter Selections**
- Conclusion

Tradeoff: Timer versus Counter

Tradeoff: Timer versus Counter

Low load: λ < $N-1$ τ $\overline{V} \approx \frac{1}{2}$ $\frac{1}{\lambda} + \tau + T_w$

ii High load:
$$
\lambda > \frac{N-1}{\tau}
$$
, $\overline{V} \approx \frac{N}{\lambda} + T_w$

Medium load: $\lambda \approx$ $N-1$ τ $,\bar{V}\geq\frac{1}{2}$ $\frac{1}{\lambda} + \tau + T_w \approx$ \overline{N} $\frac{N}{\lambda}+T_w$

Approximation of Mean Vacation Time

$$
\bar{V} \approx min\{\bar{V}_{\tau}, \bar{V}_{N}\} = min\left\{\frac{1}{\lambda} + \tau, \frac{N}{\lambda}\right\} + T_{W}
$$

Approximation of Performances

$$
\eta \approx \min\{\eta_{\tau}, \eta_{N}\} = \min\left\{1 - \frac{T_{S} + T_{W}}{\frac{1}{\lambda} + \tau + T_{W}}, 1 - \frac{T_{S} + T_{W}}{\frac{N}{\lambda} + T_{W}}\right\} \times \frac{(1 - \rho) \times (\varphi_{h} - \varphi_{l})}{\varphi_{h}}
$$

$$
D \approx \min\{D_{\tau}, D_{N}\} = \min\left\{\frac{(\lambda \tau + \lambda T_{W})^{2} + 2(\lambda \tau + \lambda T_{W})}{2\lambda(1 + \lambda \tau + \lambda T_{W})}, \frac{(N + \lambda T_{W})^{2} - N}{2\lambda(N + \lambda T_{W})}\right\} + \frac{\lambda \overline{X^{2}}}{2(1 - \rho)} + \overline{X}
$$

Optimal Relation of Timer and Counter

 τ &N policy can adapts to traffic fluctuations and avoids large delays, especially in two side regions.

Rule EEE 1

EEE 1: For a given steady state traffic rate λ , the selection of parameters τ and N should comply with the following condition:

$$
\frac{N-1}{\tau}=\lambda.
$$

Power Efficiency versus Mean Delay

- Excessive large τ and N degrade delay performance while marginally enhancing the power efficiency.
	- **a** (a) With the increase of τ , N , $\bar{V} \rightarrow \infty$ and $\eta \rightarrow \eta^*$
	- (b) D is almost linearly proportional to τ and N

Mean delay: a function of the power efficiency

$$
D \approx \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{T_s + T_w}{2\left(1 - \frac{\eta}{\eta^*}\right)} - \frac{T_s + T_w \frac{\eta}{\eta^*}}{2\lambda (T_s + T_w)} + \overline{X},
$$

The derivative

$$
\frac{dD}{d\eta} \approx \frac{T_s + T_w}{2\eta^* \left(1 - \frac{\eta}{\eta^*}\right)^2} - \frac{T_w}{2\lambda\eta^*(T_s + T_w)}
$$

EEE 2: Parameter *N* of the EEE protocol can be selected according to a given average delay requirement D from the expression of D_N .

$$
D_N \approx \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{(N + \lambda T_W)^2 - N}{2\lambda (N + \lambda T_W)} + \overline{X}
$$

Conclusions

- Develop a new approach to analyze the $M/G/1$ queue with the vacation time that is governed by the arrival process and the parameters τ and N.
- Derive a generalized P-K formula of mean delay

$$
D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{H''(1)}{2\lambda H'(1)} + \overline{X}
$$

- **Provide two rules to select appropriate** τ **and N.**
	- EEE 1

$$
\frac{N-1}{\tau} = \lambda
$$

EEE 2

$$
D_N \approx \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{(N + \lambda T_W)^2 - N}{2\lambda (N + \lambda T_W)} + \overline{X}
$$

Thanks for your attention

Our work has been uploaded to arXiv.org and is available at https://arxiv.org/abs/1611.04394