

Power Efficiency and Delay Tradeoff of Energy Efficient Ethernet Protocol

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Outline



Background

- Energy Efficient Ethernet Protocol
- Vacation Model
- Power Efficiency
- P-K Formula of Mean Delay
- Tradeoff and Parameter Selections
- Conclusion

Widely Applied Ethernets





Data Center^[1]











A. Greenberg, J.R. Hamilton, N. Jain, et al, "VL2: a scalable and flexible data center network," *Proc. ACM SIGCOMM*, 2009, pp. 51-62.
 M. Huynh, P. Mohapatra, "Metropolitan Ethernet Network: A move from LAN to MAN," *Computer Networks*, vol. 51, pp. 4867-4894, Dec 2007.
 A. Kasim, P. Adhikari, N. Chen, et al, "Ethernet: From LAN to the WAN," in *Delivering Carrier Ethernet*, 1st ed., New York: McGraw-Hill, 2007, pp.3-43.

Growing of Ethernet Devices



• The number of devices is huge and still grows rapidly.



Network Equipment Market Scale and Forecast of China

[1] http://www.ccwresearch.com.cn/view_point_detail.htm?id=557063

[2] R. Bolla, R. Bruschi, F. Davoli, and F. Cucchietti, "Energy efficiency in the future Internet: A survey of existing approaches and trends in energyaware fixed network infrastructures," *IEEE Communications Surveys Tutorials*, vol. 13, pp. 223–244, Second 2011.

Increase of Data Rate



10000



P. J. Winzer, "Beyond 100G Ethernet," *IEEE Communications Magazine*, vol. 48, pp. 26–30, July 2010.
 P. Reviriego, K. Christensen, J. Rabanillo, and J. A. Maestro, "An initial evaluation of Energy Efficient Ethernet," *IEEE Communications Letters*, vol. 15, pp. 578–580, May 2011.
 B. Kohl, "10GBASE-T power budget summary," 2007.

Idea of Energy Saving



 IEEE 802.3az: Shut down some component during idle periods and make the system more energy proportional to load



Trace from LBNL: File Server with 1G Ethernet Link

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Energy Efficiency Ethernet Protocol



- Sleep: transition time from Active to LPI
- Wakeup: transition time from LPI to Active
- LPI: low power idle mode
- Active: packets transmission period



A Typical State Transition and Power Consumption of EEE Protocol



- Counter *N*
 - Bound the backlogged queue length
- Timer τ ($\tau > T_s$)
 - Bound the delay





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 $\tau \& N$ policy

τ policy and N policy







Performance Tradeoff



- Power efficiency is improved at the expanse of delay.
- How to select N and τ to optimize system performances?



Our Works



- Model BTR strategy as an M/G/1 queue with vacation time which is governed by the arrival process.
- Derive the P-K formula of mean delay.
- Demonstrate the impacts of counter and timer on performances and provide two rules to select appropriate parameters N and τ .

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Cycles of EEE Working Process

• A renewal cycle (C)=Vacation period(V)+Busy period(B)



EEE Working Process

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EEE Working Process



Vacation period depends on the probability h_n
 h_n = Pr{n arrivals during a vacation period V}





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Six mutually independent events





 $a_n (0 < n < N)$





$$a_{n} (n = N)$$

$$a_{n} (n = N$$

Probability h_n



$$a_{n} = \begin{cases} 0, & n = 0\\ e^{-\lambda \tau} \frac{(\lambda \tau)^{n-1}}{(n-1)!}, & n = 1, 2, \cdots, N-1\\ \sum_{n=0}^{N} e^{-\lambda T_{s}} \frac{(\lambda T_{s})^{n}}{n!} - \sum_{n=1}^{N-1} e^{-\lambda \tau} \frac{(\lambda \tau)^{n-1}}{(n-1)!}, & n = N\\ e^{-\lambda T_{s}} \frac{(\lambda T_{s})^{n}}{n!}, & n = N+1, N+2, \cdots \end{cases}$$

 $b_n = Pr\{n \text{ arrivals during } T_w\} = e^{-\lambda T_w \frac{(\lambda T_w)^n}{n!}}$

 $\rightarrow h_n = \sum_{k=0}^n a_{n-k} b_k \rightarrow H(z) = A(z)B(z)$



- Mean number of arrivals during vacation is $\bar{\alpha} = H'(1)$
 - $H(z) = \sum_{n=0}^{\infty} h_n z^n$
 - $\bar{\alpha}$: mean number of arrivals during vacation period
- By Little's Law, the mean vacation time \overline{V} is

$$\bar{\alpha} = \lambda \bar{V} \to \bar{V} = \frac{\bar{\alpha}}{\lambda}$$

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Power Efficiency η

$$\eta = \frac{\varphi_h - \varphi_{EEE}}{\varphi_h} = \left(1 - \frac{T_w + T_s}{\overline{V}}\right) \cdot \frac{(1 - \rho) \times (\varphi_h - \varphi_l)}{\varphi_h}$$

When *N* and τ come to infinite

$$\eta^* = \lim_{\overline{V} \to \infty} \eta = \frac{(1-\rho) \times (\varphi_h - \varphi_l)}{\varphi_h}$$



States

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Classical M/G/1 with Vacation System



• Vacation time distribution is independent of the arrival process $H(z) = \sum_{n=0}^{\infty} h_n z^n = \sum_{n=0}^{\infty} \int_0^{\infty} (\lambda x)^n \frac{1}{n!} e^{-\lambda x} dV(x) z^n$ $= V^* (\lambda - \lambda z)$



Classical M/G/1 with Vacation

Failure of
$$H(z) = V^*(\lambda - \lambda z)$$



 In EEE protocol, the vacation time is completely governed by the arrival process. Take τ policy for example

$$V = I_0 + \tau + T_w$$

$$v(x) = \lambda e^{-\lambda(x-\tau-T_W)} (x \ge \tau + T_W) \to V^*(s) = \frac{\lambda}{\lambda+s} e^{-s(\tau+T_W)}$$

$$h_n = e^{-\lambda(\tau + T_w)} \frac{[\lambda(\tau + T_w)]^{n-1}}{(n-1)!} (n \ge 1) \to H(z) = z e^{-\lambda(1-z)(\tau + T_w)}$$

• Obviously, $H(z) \neq V^*(\lambda - \lambda z)$





• P-K Formula in the classical M/G/1 queue with vacation time

$$D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{\overline{V^2}}{2\overline{V}} + \overline{X}$$

For the same reason, it fails in EEE systems





• The waiting time for a frame is constituted by three parts.



(b) Frame F_i arrive during busy period





• *R*: mean residual time

Mean Delay





$R = \boldsymbol{E}[\boldsymbol{R}_i | \boldsymbol{\xi} = \boldsymbol{0}] \times Pr\{\boldsymbol{\xi} = \boldsymbol{0}\} + \boldsymbol{E}[\boldsymbol{R}_i | \boldsymbol{\xi} = \boldsymbol{1}] \times Pr\{\boldsymbol{\xi} = \boldsymbol{1}\}$ $= E[\boldsymbol{R}_i | \boldsymbol{\xi} = \boldsymbol{0}] \times (\boldsymbol{1} - \boldsymbol{\rho}) + E[\boldsymbol{R}_i | \boldsymbol{\xi} = \boldsymbol{1}] \times \boldsymbol{\rho}$

 $\xi = \begin{cases} 0, & \text{if a arrival comes during a vacation period} \\ 1, & \text{if a arrival comes during a busy period} \end{cases}$

$$E[R_i|\xi = 1] = \frac{1}{2\rho}\lambda \overline{X^2}^{[1]}$$
$$E[R_i|\xi = 0] = ?$$

Residual Vacation Time of Each Arrival



• When given V_n , # of arrival during the residual vacation time seen by a frame is determined.





 $E[R_i|\xi = 0] = \sum_{n=1}^{\infty} E[R_i|\xi = 0, frame \ i \ arrives \ in \ a \ V_n] \cdot P_n$

Applying Little's Law

 $\lambda E[R_i|\xi = 0] = \sum_{n=1}^{\infty} \lambda E[R_i|\xi = 0, frame \ i \ arrives \ in \ a \ V_n] \cdot P_n$

 $=\sum_{n=1}^{\infty} E[Q_i|\xi=0, frame \ i \ arrives \ V_n] \cdot P_n$

$$= \sum_{n=1}^{\infty} \left[\frac{(n-1)+(n-2)+\dots+1+0}{n} \right] \cdot \frac{n \cdot h_n}{H'(1)} = \frac{H''(1)}{2H'(1)}.$$



• **Theorem 1:** The mean delay of EEE systems is given by:

$$D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{H''(1)}{2\lambda H'(1)} + \overline{X}$$

Classical P-K Formula

$$D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{\overline{V^2}}{2\overline{V}} + \overline{X}$$

• When $H(z) = V^*(\lambda - \lambda z)$ holds, $D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{H''(1)}{2\lambda H'(1)} + \overline{X}$ degenerate into the classical form.

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Tradeoff: Timer versus Counter





Tradeoff: Timer versus Counter



• Low load: $\lambda < \frac{N-1}{\tau}, \ \overline{V} \approx \frac{1}{\lambda} + \tau + T_w$

• High load:
$$\lambda > \frac{N-1}{\tau}, \overline{V} \approx \frac{N}{\lambda} + T_w$$

• Medium load: $\lambda \approx \frac{N-1}{\tau}, \overline{V} \ge \frac{1}{\lambda} + \tau + T_w \approx \frac{N}{\lambda} + T_w$



Approximation of Mean Vacation Time



$$\overline{V} \approx \min\{\overline{V}_{\tau}, \overline{V}_{N}\} = \min\{\frac{1}{\lambda} + \tau, \frac{N}{\lambda}\} + T_{w}$$



Approximation of Performances



$$\eta \approx \min\{\eta_{\tau}, \eta_{N}\} = \min\left\{1 - \frac{T_{S} + T_{W}}{\frac{1}{\lambda} + \tau + T_{W}}, 1 - \frac{T_{S} + T_{W}}{\frac{N}{\lambda} + T_{W}}\right\} \times \frac{(1 - \rho) \times (\varphi_{h} - \varphi_{l})}{\varphi_{h}}$$
$$D \approx \min\{D_{\tau}, D_{N}\} = \min\left\{\frac{(\lambda \tau + \lambda T_{W})^{2} + 2(\lambda \tau + \lambda T_{W})}{2\lambda(1 + \lambda \tau + \lambda T_{W})}, \frac{(N + \lambda T_{W})^{2} - N}{2\lambda(N + \lambda T_{W})}\right\} + \frac{\lambda \overline{X^{2}}}{2(1 - \rho)} + \overline{X}$$



Optimal Relation of Timer and Counter



τ&*N* policy can adapts to traffic fluctuations and avoids large delays, especially in two side regions.



Rule EEE 1



EEE 1: For a given steady state traffic rate λ , the selection of parameters τ and *N* should comply with the following condition:

$$\frac{N-1}{\tau} = \lambda.$$

Power Efficiency versus Mean Delay



- Excessive large τ and N degrade delay performance while marginally enhancing the power efficiency.
 - (a) With the increase of τ , N, $\overline{V} \to \infty$ and $\eta \to \eta^*$
 - (b) D is almost linearly proportional to τ and N





Mean delay: a function of the power efficiency

$$D \approx \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{T_s + T_w}{2\left(1 - \frac{\eta}{\eta^*}\right)} - \frac{T_s + T_w \frac{\eta}{\eta^*}}{2\lambda(T_s + T_w)} + \overline{X},$$

The derivative

$$\frac{dD}{d\eta} \approx \frac{T_S + T_W}{2\eta^* \left(1 - \frac{\eta}{\eta^*}\right)^2} - \frac{T_W}{2\lambda \eta^* (T_S + T_W)}$$





EEE 2: Parameter *N* of the EEE protocol can be selected according to a given average delay requirement *D* from the expression of D_N .

$$D_N \approx \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{(N+\lambda T_w)^2 - N}{2\lambda(N+\lambda T_w)} + \overline{X}$$

Conclusions



- Develop a new approach to analyze the *M/G/1* queue with the vacation time that is governed by the arrival process and the parameters *τ* and *N*.
- Derive a generalized P-K formula of mean delay

$$D = \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{H^{\prime\prime}(1)}{2\lambda H^{\prime}(1)} + \overline{X}$$

- Provide two rules to select appropriate τ and N.
 - **EEE** 1

$$\frac{N-1}{\tau} = \lambda$$

• EEE 2

$$D_N \approx \frac{\lambda \overline{X^2}}{2(1-\rho)} + \frac{(N+\lambda T_w)^2 - N}{2\lambda (N+\lambda T_w)} + \overline{X}$$



Thanks for your attention

Our work has been uploaded to arXiv.org and is available at https://arxiv.org/abs/1611.04394